

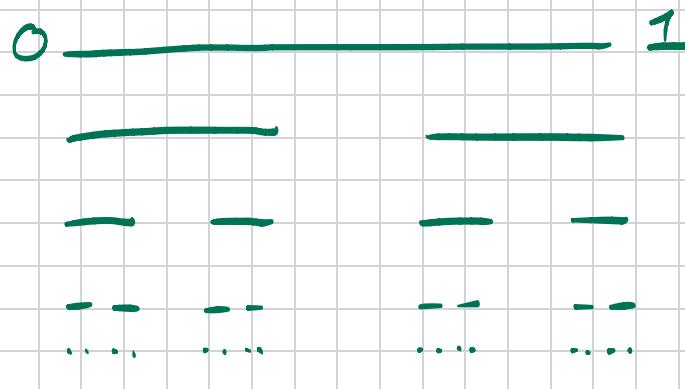


Introduction to Fractal Geometry

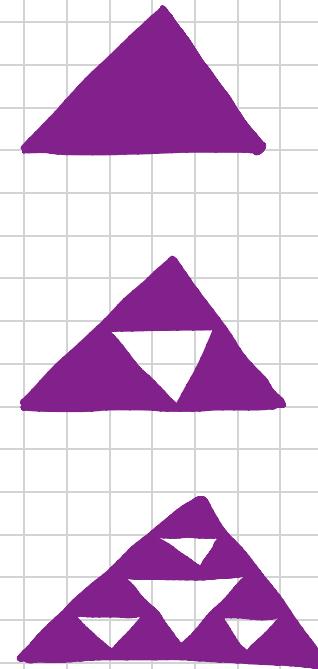
SUMS 2024

GMS

I Fractal Objects



Middle-third Cantor Set



Sierpinski triangle

Mandelbrot Set



colour = rate at which
 $Z_n \rightarrow +\infty$, where
 $Z_0 = 0, Z_{n+1} = Z_n^2 + c$

② Fractal Experts



Benoit B.
Mandelbrot

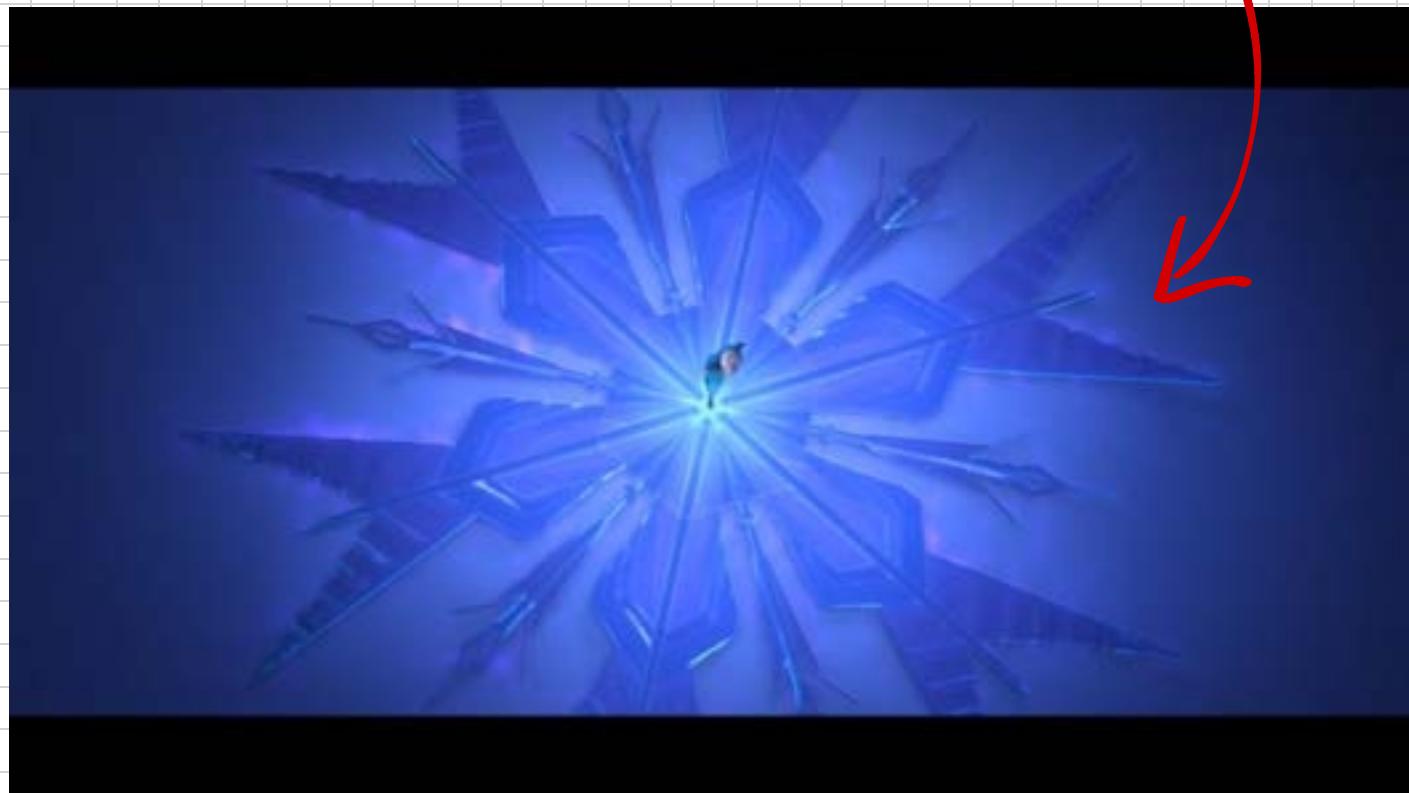


Prof Kenneth Falconer
(Regius Professor of
Mathematics, St.A)

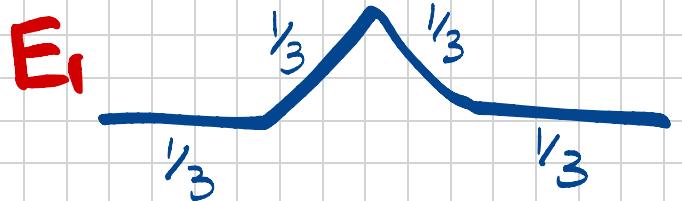


Queen
Elsa

"Let it go" — fractal snowflakes



③ koch curve



Define the fractal as $F = \bigcap_{k=0}^{\infty} E_k$

Let $L(k)$ be the length of E_k , and $A(k)$ the area enclosed by it and $y=0$.

$$L(0) = 1, \quad A(0) = 0$$

$$L(1) = \frac{4}{3}, \quad A(1) = \frac{\sqrt{3}}{36}$$

$$L(2) = \left(\frac{4}{3}\right)^2, \quad A(2) = \frac{\sqrt{3}}{36} \left(1 + \frac{4}{9}\right)$$

⋮

$$L(k) = \left(\frac{4}{3}\right)^k \rightarrow +\infty$$

$$A(k) = \frac{\sqrt{3}}{36} \sum_{i=0}^k \left(\frac{4}{9}\right)^i$$

$$\rightarrow \frac{\sqrt{3}}{36} \left(\frac{1}{1 - \frac{4}{9}}\right) = \frac{\sqrt{3}}{20}$$

Area of is $\frac{\sqrt{3}}{4} l^2$

④

Natural fractals



self-
similar



How long
is this
coast?
(Will
return to
this!)



What is the
Scale?



↑
fine
detail

branching
structure

⑤ What is a fractal?

all of them!

and
scale-
invariance

- fine structure / detail at multiple scales
- too "irregular" for length, area etc. to satisfy
- Self-similarity (small parts look like the whole)
- usually also simple in some way (e.g. recursive definition)
- "fractal dimension" greater than topological dimension

Koch curve

Fern

Cantor set,
Koch curve

0 = disconnected points

1 = line

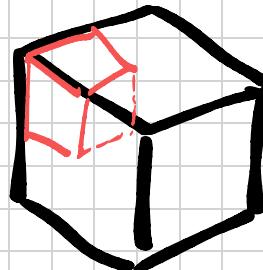
2 = area etc.

(Falconer,
"Fractal Geometry"
2014)

⑥

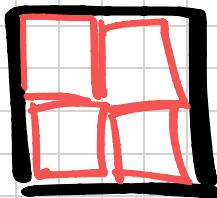
Dimension

$$N \sim \delta^{-d}$$



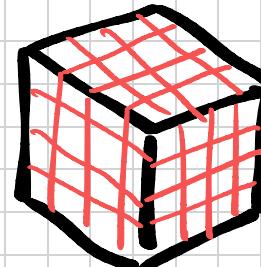
$$\delta = \frac{1}{2}$$
$$N = 8$$

$$\delta = \frac{1}{2}$$
$$N = 2$$



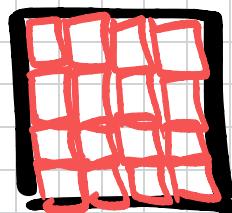
$$\delta = \frac{1}{2}$$
$$N = 4$$

$$\delta = \frac{1}{4}$$
$$N = 16$$

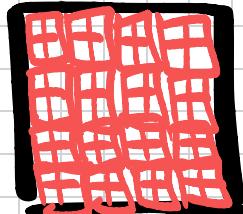


$$\delta = \frac{1}{4}$$
$$N = 64$$

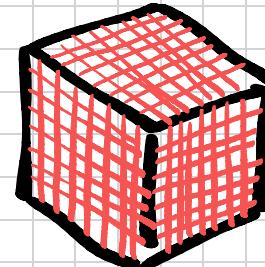
$$\delta = \frac{1}{4}$$
$$N = 4$$



$$\delta = \frac{1}{8}$$
$$N = 8$$



$$\delta = \frac{1}{8}$$
$$N = 64$$



$$\delta = \frac{1}{8}$$
$$N = 512$$

$d = 1$

$d = 2$

$d = 3$

⑦ Box-counting Dimension

Let $F \subseteq \mathbb{R}^n$ be non-empty and bounded.

For $\delta > 0$, $\{U_i\}_{i=1}^m$ is a δ -cover of F if:

$$F \subseteq \bigcup_i U_i \quad \text{and} \quad 0 \leq |U_i| \leq \delta \quad \forall i$$

(they cover F)

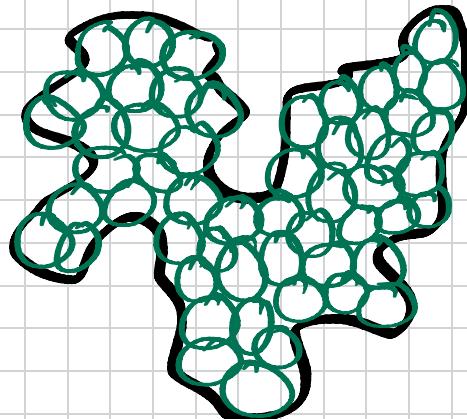
(they are each δ -wide)

Let $N_\delta(F)$ be the least number of sets in a δ -cover of F .

Then

$$\dim_B(F) = \lim_{\delta \rightarrow 0} \frac{\log(N_\delta(F))}{\log(\delta^{-1})}$$

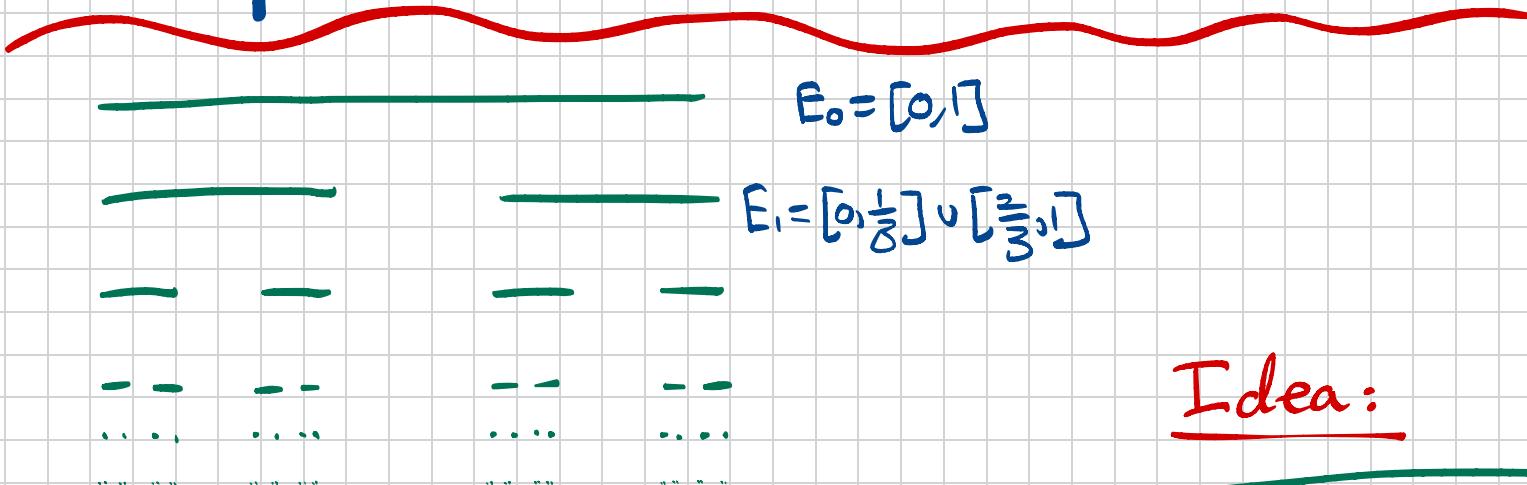
"How many boxes of size δ do we need to cover the set?"



if this limit exists.

↙ diameter $\leq \delta$

Example : Middle-third Cantor Set



Level E_k consists
of 2^k intervals of length 3^{-k}

$$F = \bigcap_{n=0}^{\infty} E_n \subseteq E_h \text{ for any } h \in \mathbb{N}_0$$

Idea:

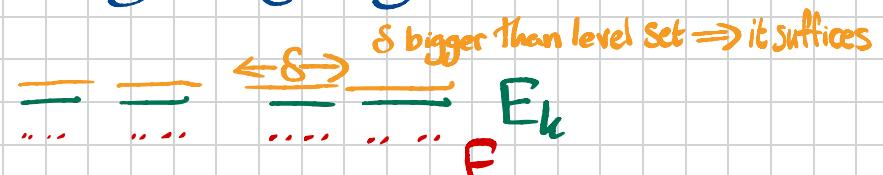
Obtain upper and lower bounds
for the minimum δ -cover at a
given level.

Then take the limit as $\delta \rightarrow 0$

Upper bound: ("How many sets would suffice?")

Let $0 < \delta < 1$, then \exists unique k s.t.

$$3^{-k} < \delta \leq 3^{-k+1}$$



So E_k is a δ -cover of F , with 2^k sets. Hence $N_\delta(F) \leq 2^k$

$$\therefore \frac{\log(N_\delta(F))}{-\log(\delta)} \leq \frac{\log(2^k)}{-\log(3^{-k+1})}$$

$$= \frac{k \log(2)}{(k-1) \log(3)} \rightarrow \frac{\log(2)}{\log(3)} \text{ as } \delta \rightarrow 0 \quad (\Rightarrow k \rightarrow +\infty)$$

Lower bound: ("How many MUST I need?")

Let $0 < \delta < 1$, then \exists unique k s.t.

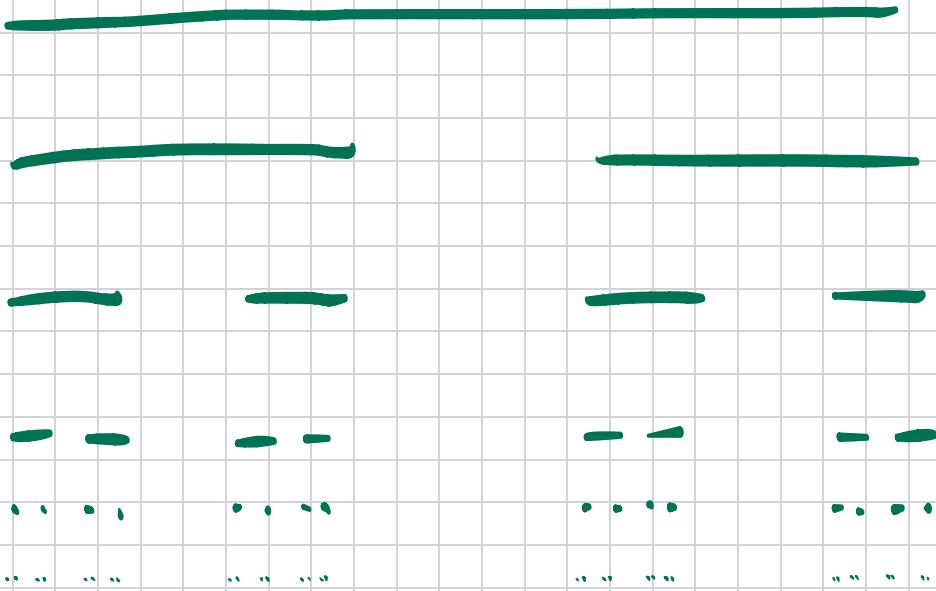
$$3^{-k+1} \leq \delta < 3^{-k}$$



Since δ is smaller than the gaps, any set in a δ -cover intersects only one interval in E_k — So at least one is required for each: $N_\delta(F) \geq 2^k$

$$\therefore \frac{\log(N_\delta(F))}{-\log(\delta)} \geq \frac{\log(2^k)}{-\log(3^{-k+1})} = \frac{k \log(2)}{(k-1) \log(3)}$$
$$\rightarrow \frac{\log(2)}{\log(3)} \text{ as } \delta \rightarrow 0.$$

Hence in the limit $\delta \rightarrow 0$,



$$\frac{\log(2)}{\log(3)} \leq \dim_B(F) \leq \frac{\log(2)}{\log(3)}$$

So the box-counting dimension is:

$$\dim_B(F) = \frac{\log(2)}{\log(3)} \approx 0.631$$

II Divider dimension

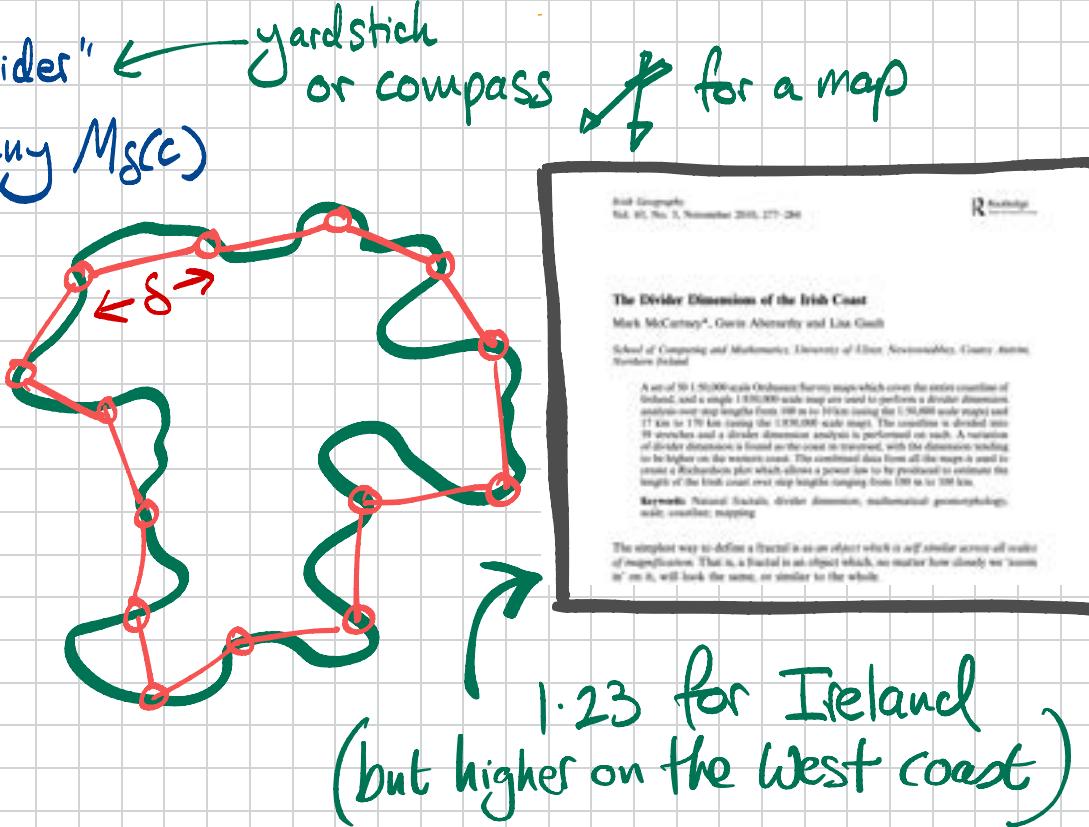
(There are many related definitions of "dimension"!)

- Given a curve C , set "divider" width S and measure how many $M_S(C)$ points S -apart on the curve.

- Repeat for smaller S .

- Estimate log-log gradient:

$$\dim_D(C) = \lim_{S \rightarrow 0} \frac{\log(M_S(C))}{\log(S^{-1})}$$

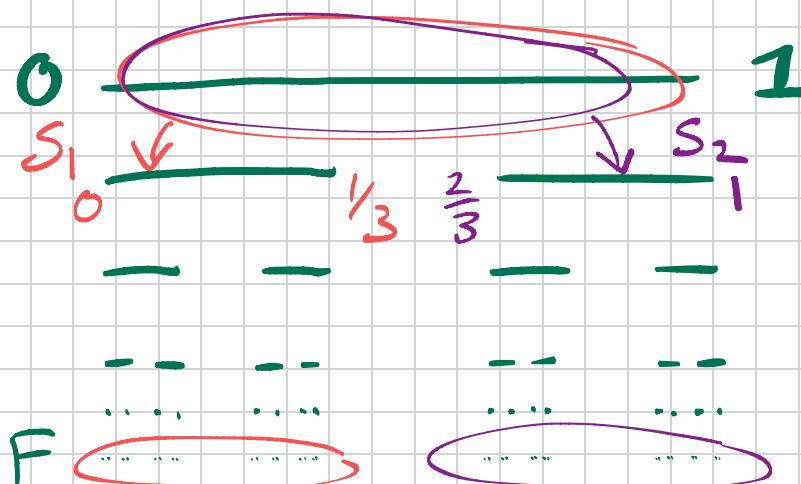


(12)

Iterated Function Systems

Alternatively, the non-empty compact invariant set associated with a system of contracting functions:

$$F = \bigcup_{i=1}^N S_i(F)$$



Two copies of the whole set.

For middle-third
Cantor set:

$$S_1(x) = \frac{1}{3}x, \quad S_2(x) = \frac{1}{3}x + \frac{2}{3}$$

Then F is uniquely such that:

$$F = S_1(F) \cup S_2(F)$$

Let

$S: \mathbb{R}^n \rightarrow \mathbb{R}^n$ be

$$S(x) = \bigcup_{i=1}^N S_i(x)$$

Then for suitable
 $A \subseteq \mathbb{R}^n$,

$$F = \overline{\bigcap_{k=1}^{\infty} S^k(A)}$$

So choosing a suitable system of contractions
 S_1, \dots, S_N and starting sets $A \subseteq \mathbb{R}^2$ the
system converges to a fractal.

↑
generate your own at:

github.com/gavinabernethy/gifs_generator

References

key
textbook

} Falconer, K. (2014) Fractal geometry: mathematical foundations and applications.
John Wiley & Sons.

McCartney, M. (2020). The area, centroid and volume of revolution of the Koch curve.
International Journal of Mathematical Education in Science and Technology, 52(5), 782–786.